

Rooted Trees

CMP-225




Graphs

Graph: a pair $G = (V, E)$, with

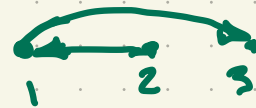
- V a set called "vertices" or "nodes"

- E a set of pairs from V , i.e. $E \subseteq V \times V$, called edges.

Eg: $G = (\{1, 2, 3\}, \{(1, 2), (1, 3)\})$  (*)

" G is directed": edges are ordered pairs (often called "arcs"):

$\langle \{1, 2, 3\}, \{(1, 2), (1, 3)\} \rangle \neq \langle \{1, 2, 3\}, \{(2, 1), (1, 3)\} \rangle$



" G is undirected": edges are sets

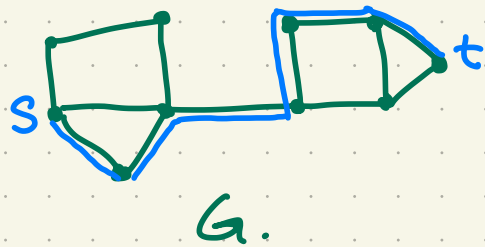
$\langle \{1, 2, 3\}, \{(2, 1), (1, 3)\} \rangle = \langle \{1, 2, 3\}, \{(1, 2), (3, 1)\} \rangle = (*)$

- By default, by "graph" we will mean "undirected graph"

Path in G of length n :

- sequence $\langle v_0, v_1, v_2, \dots, v_n \rangle$ of vertices s.t.
 $(v_0, v_1), (v_1, v_2), (v_2, v_3), \dots$ are edges of G .

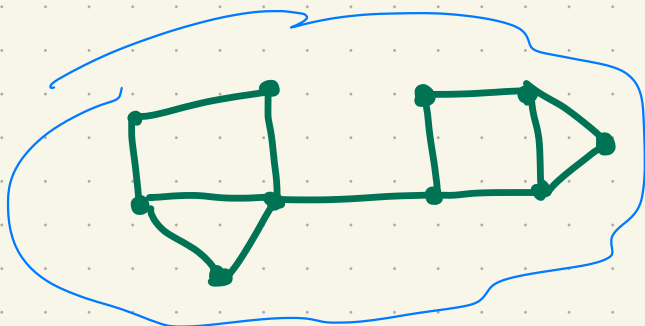
s-t Path in G : path $\langle s, \dots, t \rangle$ in G .



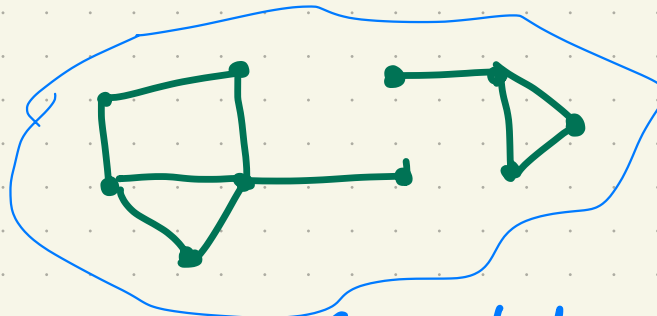
A s-t path of length 6
in G .

Vertex t is reachable from s in G if there
is an s-t path in G .

G is connected if for every pair $u, v \in V$,
 u is reachable from v .



Connected



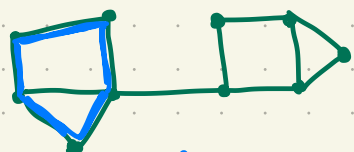
Not Connected.

Cycles & Trees

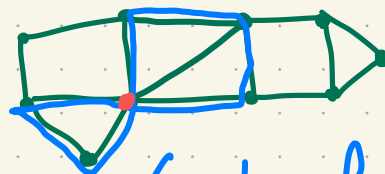
Cycle in G : path $\langle v_0, \dots, v_{n-1}, v_n \rangle$ in G where $v_0 = v_n$.

Simple Path: all vertices are distinct

Simple Cycle: cycle $\langle v_0, \dots, v_{n-1}, v_n \rangle$ where $\langle v_0, \dots, v_{n-1} \rangle$ is a simple path

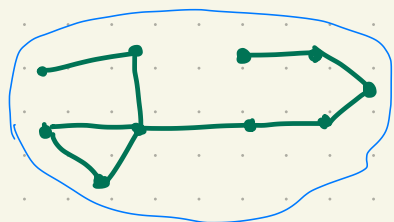


Simple cycle of length 5

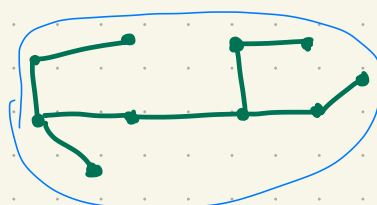


Cycle of length 7 (• repeats)

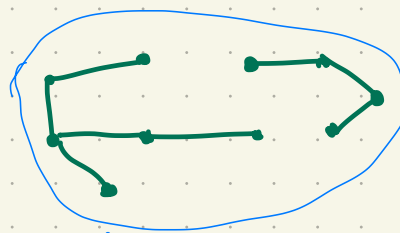
Tree: A connected, acyclic, graph.



Not a tree



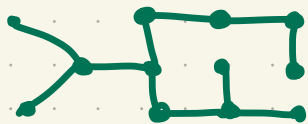
Tree



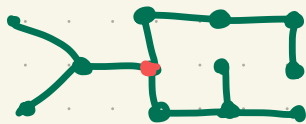
Not a tree

Fact: Every tree with n vertices has $n-1$ edges.

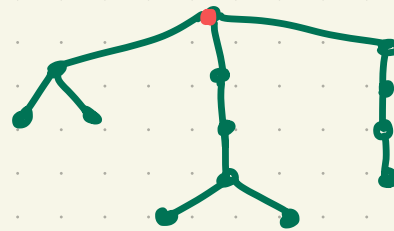
Rooted tree: tree with a distinguished vertex called the root.



Unrooted tree



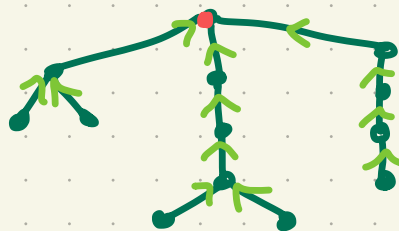
Tree T with root.



Alternate drawing of T .

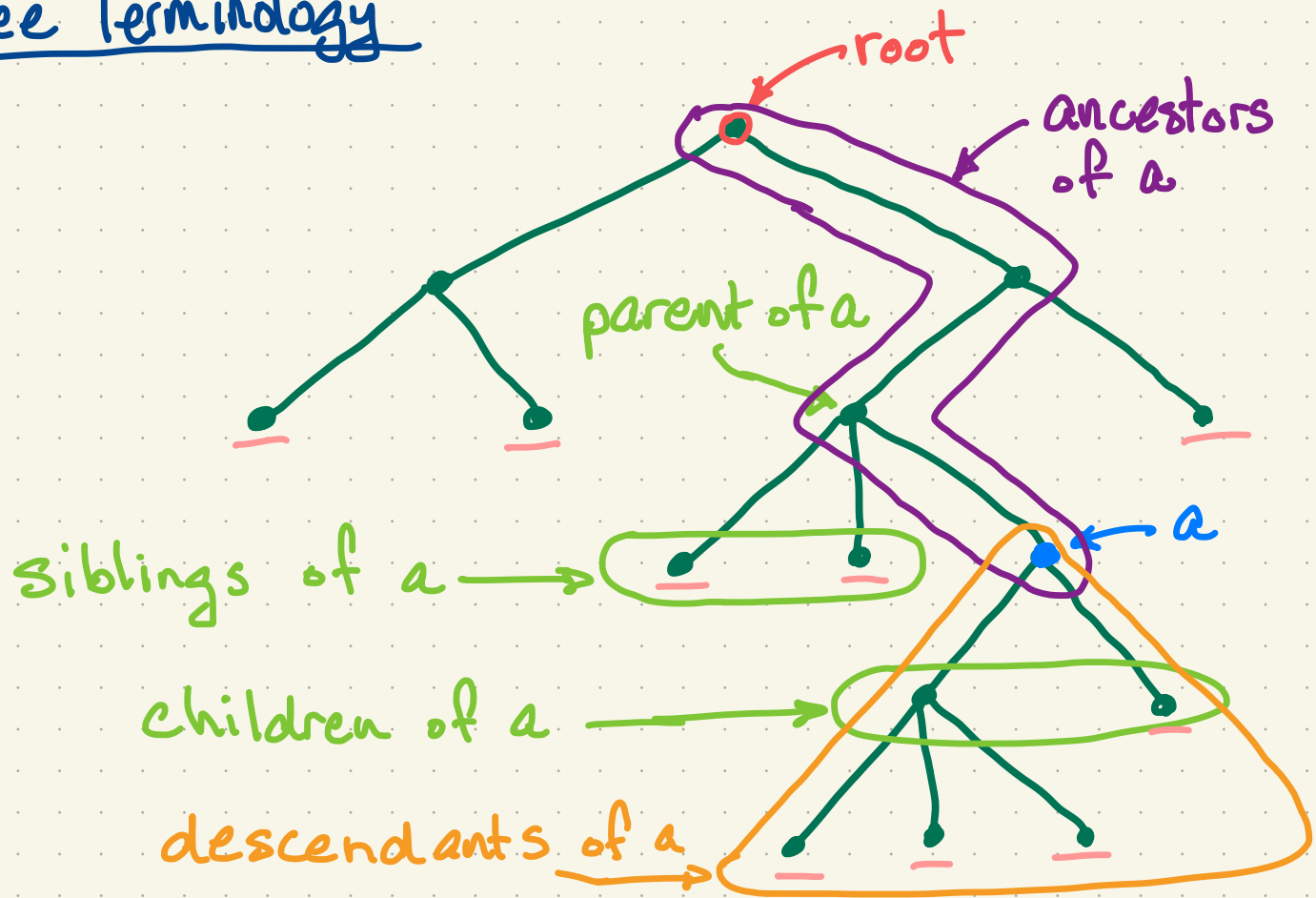
Notice: in a tree there is a unique path between any two vertices.
So: a unique path from any vertex to the root.

Thus: the root induces a direction on the edges
- eg. toward the root.



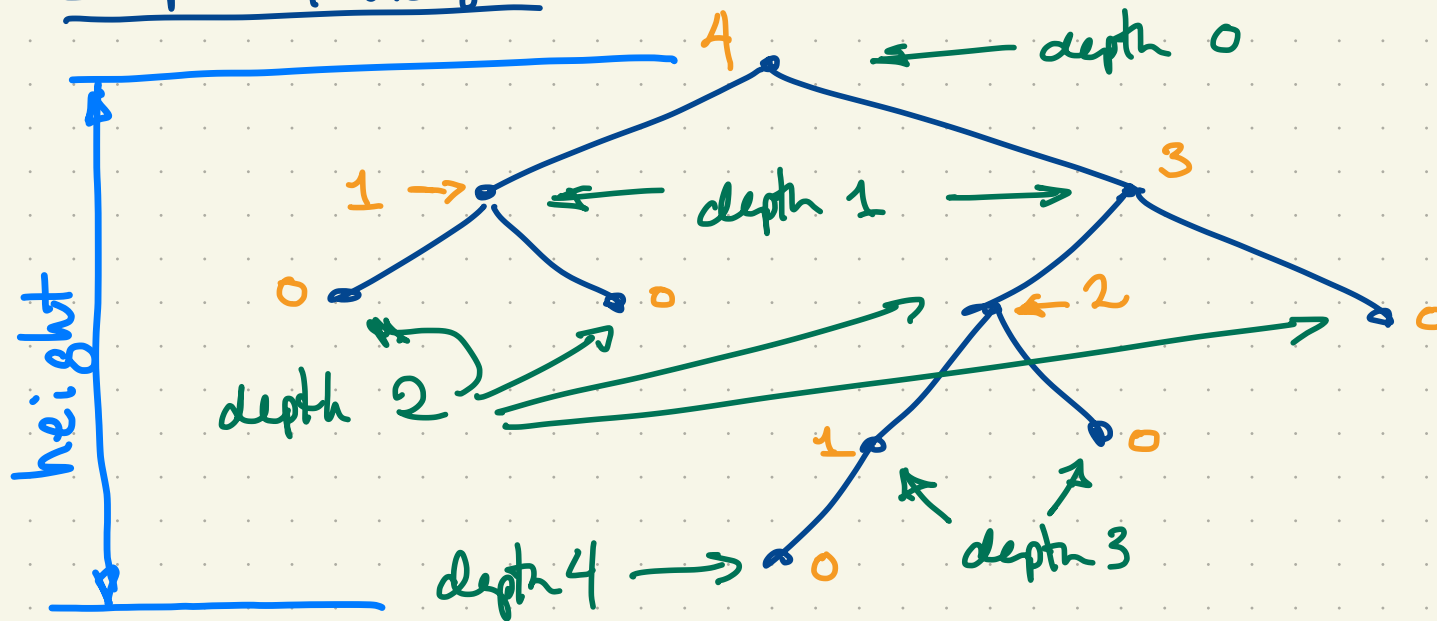
(some times "away from the root").

Rooted Tree Terminology



- The root has no parent;
- leaves have no children
- Internal nodes are the non-leaves (sometimes root excluded too).

Depth & Height



Depth of node v = length of path from v to the root.

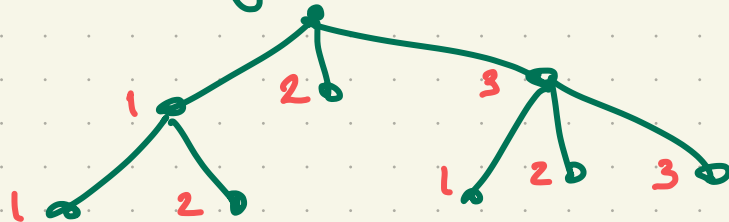
Height of node v = length of longest path from v to a descendent of v (eg. to a leaf)

Height of tree T = height of its root
= max height of any node in T
= max depth of any node in T .

A rooted tree is

- k -ary if no node has $> k$ children
- binary if no node has > 2 children
- ordered if the children of every node are ordered.

Eg: A ordered ternary tree:



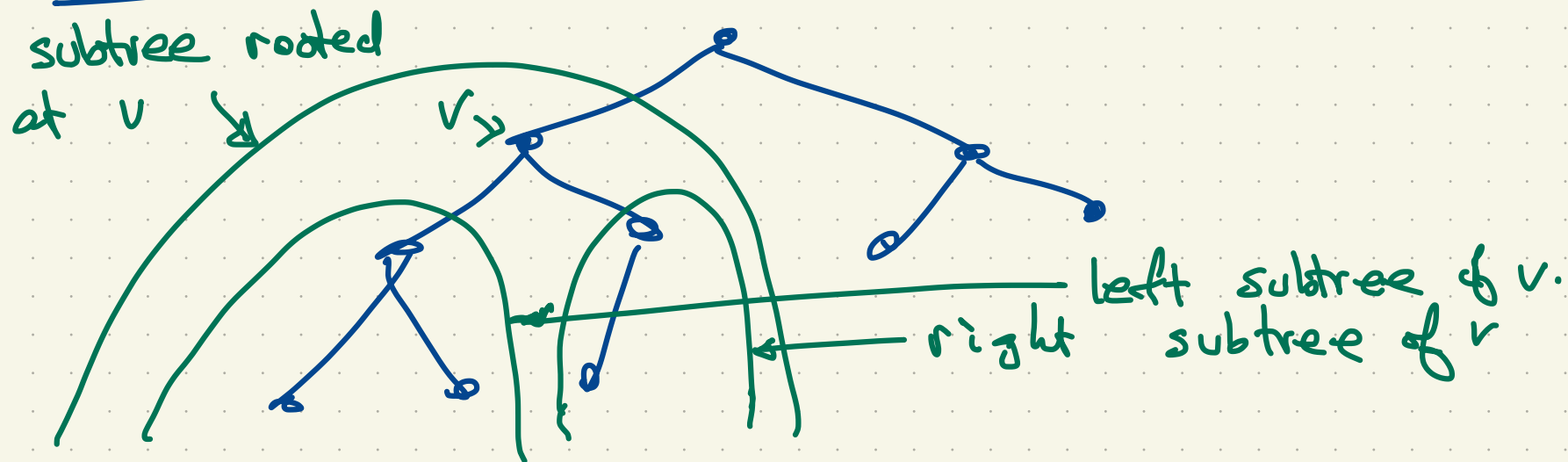
Notice: When we draw a tree, or represent it in a data structure, we order it.

In an ordered binary tree, every child of a node v is either the "left child of v " or the "right child of v ".

Subtree rooted at v : tree with root v and containing all descendants of v .

In a binary tree;

- "left subtree of v " means the subtree rooted at the left child of v
- sim. for "right subtree of v ".



End