



CMPT 295

Unit - Machine-Level Programming

Lecture 21 – Assembly language – Array – 2D

Last lecture

- Recursion
 - Handled without special instruction
 - Stack frames
 - x86-64 Function call and Register saving conventions
- Manipulation of arrays – in x86-64
 - From x86-64's perspective, an array is a contiguously allocated region of $n * L$ bytes in memory where $L = \text{sizeof}(T)$ and $T \rightarrow$ data type of elements stored in array
 - Compute memory address of each array element
 - $A[i] = A + i * L$

Today's Menu

- ▶ Introduction
 - ▶ C program -> assembly code -> machine level code
- ▶ Assembly language basics: data, `move` operation
 - ▶ Memory addressing modes
- ▶ Operation `leaq` and Arithmetic & logical operations
- ▶ Conditional Statement – Condition Code + `cmovX`
- ▶ Loops
- ▶ Function call – Stack
 - ▶ Overview of Function Call
 - ▶ Memory Layout and Stack - x86-64 instructions and registers
 - ▶ Passing control
 - ▶ Passing data – Calling Conventions
 - ▶ Managing local data
 - ▶ Recursion
- ▶ Array (cont'd)
- ▶ Buffer Overflow
- ▶ Floating-point operations

2D Array

T **A**[**R**][**C**];

=> in C

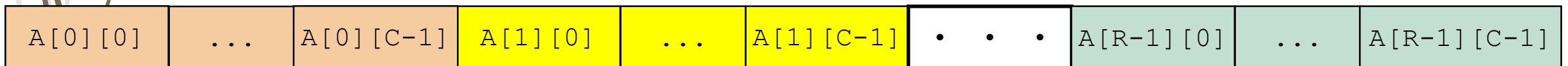
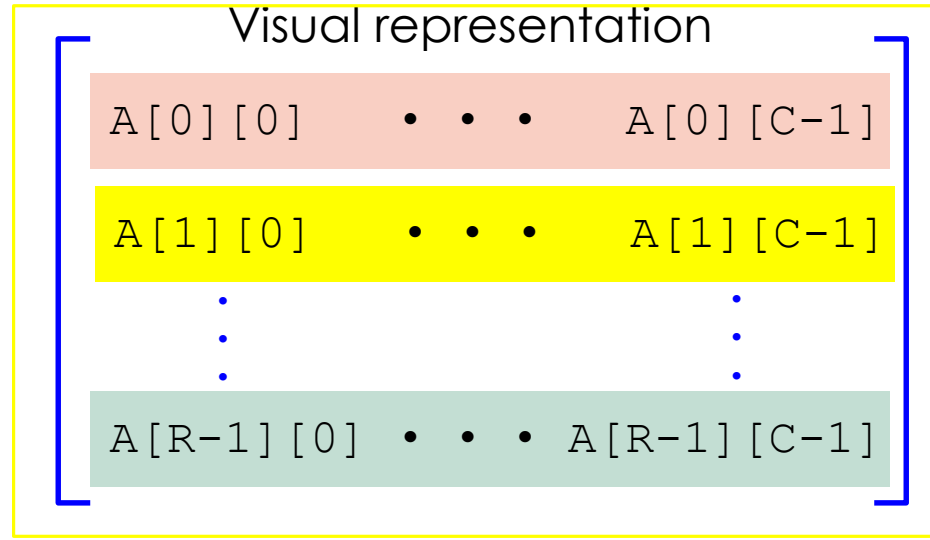
► ... is 2D array of data type **T**
R rows, **C** columns

=> in x86-64

► ... is a contiguously allocated region of **R * C * L** bytes in memory where **L = sizeof(T)**

► Array layout in memory

► **Row-Major** ordering → Example using **int A[R][C];**

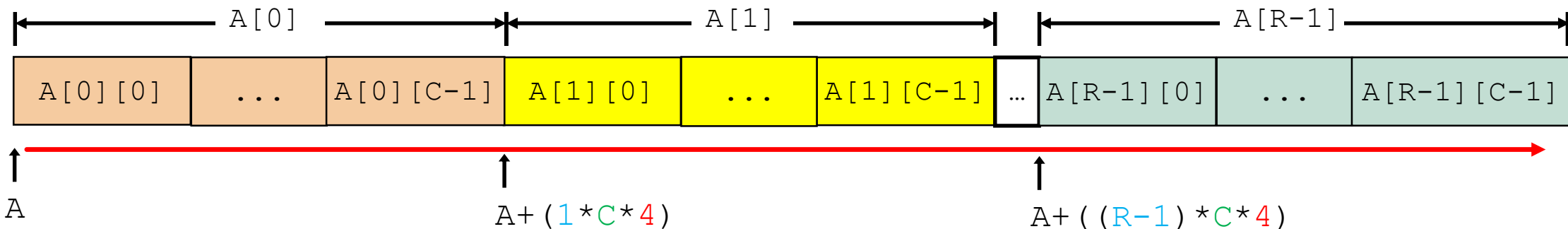


R * C * int Bytes

Accessing a row of 2D array

T **A[R][C];**

- **A[i]** is an array of **C** elements (row **i** of array **A**)
- Memory address of each row **A[i]**: $A + (i * C * L)$
 - where **A** is **base memory address**
- Can access other rows by incrementing **A** by $i * C * L$
- Example using **int A[R][C];**



Accessing an element of 2D array

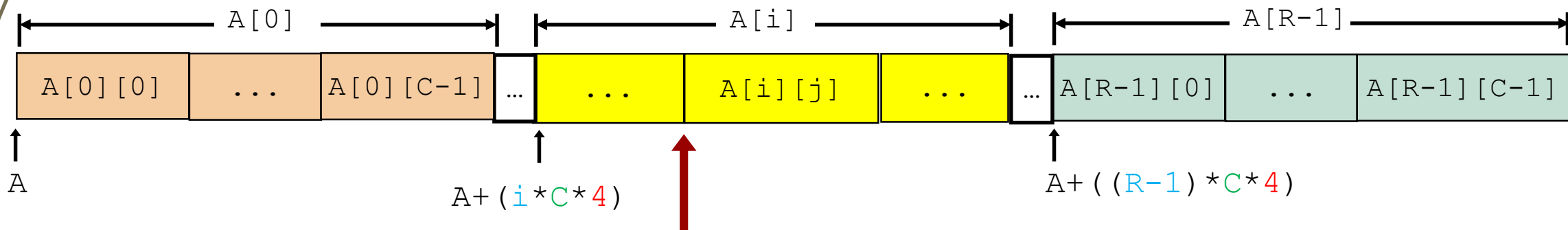
T **A**[**R**][**C**];

➤ **A**[**i**][**j**] is element of type **T**, which requires **L** bytes

➤ Memory address of each element **A**[**i**][**j**]:

$$A + (i * C * L) + (j * L) = A + (i * C + j) * L$$

➤ Example using `int A[R][C];`



Homework

Example: `int A[3][5];`

➤ In memory:



➤ Let's compute the memory address to access:

➤ `A[2]` :

➤ `A[2][3]` :

Demo - Accessing an element of 2D array

```
#define N 4

char A[N][N] = { 1, -2, 3, -4,
                -5, 6, -7, 8,
                -1, 2, -3, 4,
                5, -6, 7, -8};

char C[N][N];
```

```
void main() {

    printf("Original matrix: \n");
    printMatrixByRow(A, N);

    printf("Copy: \n");
    copy(A, C, N);
    printMatrixByRow(C, N);

    return;
}

// Print N elements in a row
void printMatrixByRow(void *D, int n) {
```

+ **copy.s** on our course web site

Summary

- Manipulation of 2D arrays – in x86-64
 - From x86-64's perspective, a 2D array is a contiguously allocated region of $R * C * L$ bytes in memory where $L = \text{sizeof}(T)$ and $T \rightarrow$ data type of elements stored in array
 - 2D Array layout in memory: Row-Major ordering
 - Memory address of each row $A[i]$: $A + (i * C * L)$
 - Memory address of each element $A[i][j]$:

$$A + (i * C * L) + (j * L)$$

$$\Rightarrow A + (i * C + j) * L$$

Next Lecture

- ▶ Introduction
 - ▶ C program -> assembly code -> machine level code
- ▶ Assembly language basics: data, `move` operation
 - ▶ Memory addressing modes
- ▶ Operation `leaq` and Arithmetic & logical operations
- ▶ Conditional Statement – Condition Code + `cmovX`
- ▶ Loops
- ▶ Function call – Stack
 - ▶ Overview of Function Call
 - ▶ Memory Layout and Stack - x86-64 instructions and registers
 - ▶ Passing control
 - ▶ Passing data – Calling Conventions
 - ▶ Managing local data
 - ▶ Recursion
- ▶ Array
- ▶ Buffer Overflow
- ▶ Floating-point operations